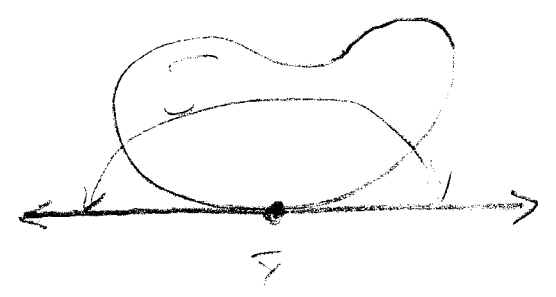
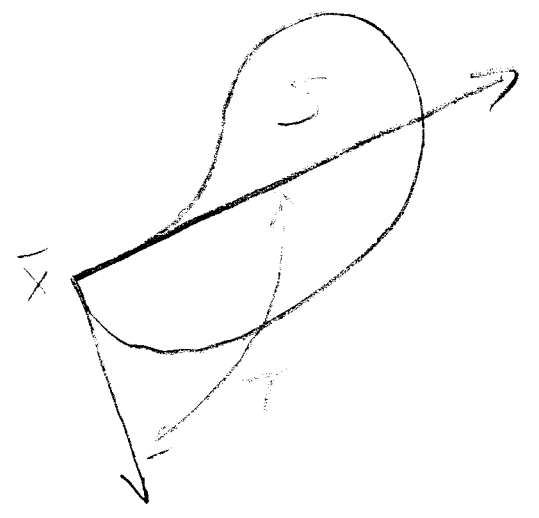


Defn $S \subseteq \mathbb{R}^n$, $\bar{x} \in \text{cl } S$.

Cone of tangents T of S at \bar{x} , is the set of all directions d such that $d = \lim_{k \rightarrow \infty} \lambda_k (x_k - \bar{x})$, for $\lambda_k > 0, x_k \in S$ for each k , and $x_k \rightarrow \bar{x}$.



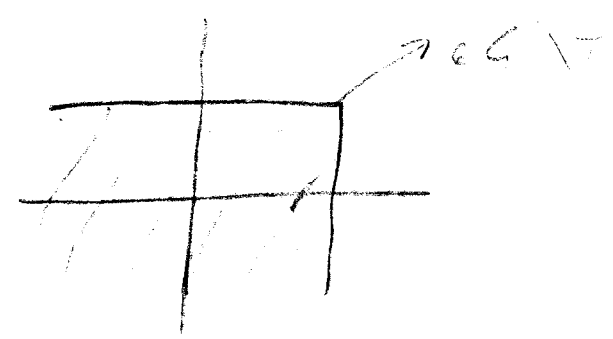
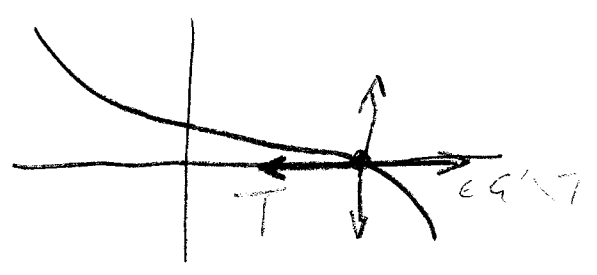
Let $G' = \{d : \nabla g_i(\bar{x})^T d \leq 0 \quad \forall i \in I\}$.

Note that $T \subseteq G'$.

Constraint Qualification:

$T = G'$

Previous examples:



Theorem If \bar{x} is a local minimum to (NLP) then $F_0 \cap T = \emptyset$,

where $F_0 = \{d : \nabla f(\bar{x})^T d < 0\}$.

Proof Let $d \in T$.

$$d = \lim_{k \rightarrow \infty} \lambda_k (x_k - \bar{x})$$

$$f(x_k) - f(\bar{x}) = \nabla f(\bar{x})^T (x_k - \bar{x}) + \|x_k - \bar{x}\| \alpha(\bar{x}, x_k - \bar{x})$$

\forall
0

So $d \notin F_0$, since $x_k \rightarrow \bar{x}$

//

Karush-Kuhn-Tucker Theorem

and CQ holds,

If \bar{x} is a local optimal solution to (NLP) then

\exists nonnegative scalars u_i for $i \in I$ such that

$$\nabla f(\bar{x}) + \sum_{i \in I} u_i \nabla g_i(\bar{x}) = 0.$$

Proof

Recall Farkas Lemma:

$$(I) \quad \begin{aligned} Ax &= b \\ x &\geq 0 \end{aligned}$$

$$(II) \quad \begin{aligned} A^T y &\leq 0 \\ b^T y &> 0 \end{aligned}$$

~~Now \bar{x} is a local optimum, so if $\nabla g_i(\bar{x})^T d \leq 0$ then $d \in C'$~~

Thus $d \in T_{C'}(\bar{x}) \Rightarrow d \in F_0$

\bar{x} is a local minimum, so $C'(\bar{x}) \cap F_0 \cap T = \emptyset$

ie. $d \notin F_0$ ie. $d^T \nabla f(\bar{x}) \geq 0$.

\therefore System (II) is infeasible \Rightarrow Farkas Lemma \Rightarrow System (I) has a solution

ie. System (I) has a solution, ie. $\exists u \geq 0$ s.t.

$$\nabla f(\bar{x}) + \sum_{i \in I} u_i \nabla g_i(\bar{x}) = 0.$$

KKT conditions in vector notation:

if \bar{x} is a local minimum ^{and CQ holds}, then $\exists \bar{u} \in \mathbb{R}^m$ with

$$\begin{aligned} \nabla f(\bar{x}) + \sum_{i=1}^m \bar{u}_i \nabla g_i(\bar{x}) &= 0 \\ \bar{u}_i g_i(\bar{x}) &= 0 \\ \bar{u}_i &\geq 0. \end{aligned}$$

Lemma ~~If $\{\nabla g_i : i \in I\}$ are lin indep, then the constraint qualification holds.~~

Lemma If $\exists \hat{d}$ s.t. $\hat{d}^T \nabla g_I < 0$ then CQ holds at $\bar{x} \in S$.

pf Let $\bar{d} \in \mathcal{D}_{S, \bar{x}}$ and set $d_k = \bar{d} + \frac{1}{k} \hat{d}$.
~~then $\nabla g_i(\bar{x})^T d_k < 0$ for all $i \in I$.~~

Lemma If $\{\nabla g_i : i \in I\}$ are lin indep, the constraint qualification holds.

Proof Consider the dual pair of LPs:

$$\begin{aligned} \max \quad & 0 \\ & \nabla g_i(\bar{x})^T d \leq -1 \quad \forall i \in I \quad (P) \\ \min \quad & -\sum w_i \\ & \sum w_i \nabla g_i(\bar{x}) = 0 \quad (D) \\ & w_i \geq 0. \end{aligned}$$

Only soln to (D) is $w = 0$.
 So $w = 0$ is the only feasible solution to (D). //

Corollary (to first lemma on p. 95)

Slater Constraint Qualification:

If $g_i(x)$ convex $\forall i$ and if $\exists \hat{x}$ with $g_i(\hat{x}) < 0 \forall i$,
then CQ holds at all feasible x .

Proof Let \bar{x} be feasible, with active set I .

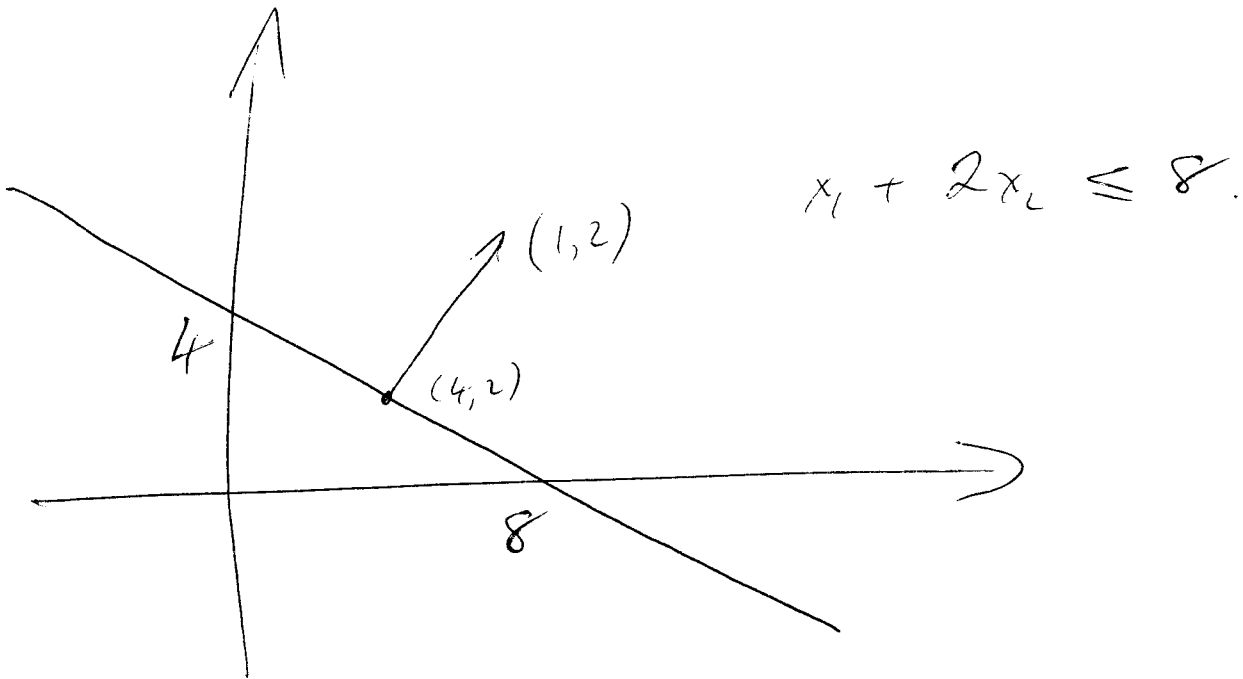
$$\text{Let } \hat{d} = \hat{x} - \bar{x}.$$

$$\text{Then } g_i(\hat{x}) < 0 = g_i(\bar{x}) \quad \forall i \in I$$

$$\text{and } g_i(\hat{x}) \geq g_i(\bar{x}) + Dg_i(\bar{x})^T(\hat{x} - \bar{x})$$

$$\text{so } Dg_i(\bar{x})^T \hat{d} < 0 \quad \forall i \in I$$





$$\begin{aligned} \min & (x_1 - 7)^2 + (x_2 - 8)^2 \\ \text{s.t.} & x_1 + 2x_2 \leq 8. \end{aligned}$$

$$\begin{aligned} \min & \frac{1}{2}x_1^2 - 7x_1 + \frac{1}{2}x_2^2 - 8x_2 \\ \text{s.t.} & x_1 + 2x_2 - 8 \leq 0 \end{aligned}$$

$$\nabla f = \begin{pmatrix} x_1 - 7 \\ x_2 - 8 \end{pmatrix}$$

$$\nabla g = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

So need solve to

$$\begin{aligned} x_1 - 7 + u &= 0 & \textcircled{1} \\ x_2 - 8 + 2u &= 0 & \textcircled{2} \end{aligned}$$

$$u(x_1 + 2x_2 - 8) = 0 \quad \textcircled{3}$$

$$x_1 + 2x_2 - 8 \leq 0 \quad \textcircled{4} \quad u_i \geq 0, \quad \textcircled{5}$$

$$u = 0 \Rightarrow x_1 = 7, x_2 = 8$$

$$\text{Violate } x_1 + 2x_2 - 8 \leq 0$$

$$u > 0 \Rightarrow x_1 + 2x_2 - 8 = 0 \quad \text{if } x_1 = 8 - 2x_2$$

Substitute into (1):

$$\begin{array}{rcl} 1 - 2x_2 + u & = & 0 \\ -8 + x_2 + 2u & = & 0 \end{array}$$

$$\therefore -15 + 5u = 0 \Rightarrow u = 3$$

$$\Rightarrow x_1 = 4, x_2 = 2. \quad \checkmark$$